

Lecture 04 – Control Flow II

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CS 343 – Fall 2020

Based on Michael Bailey's ECE 422

32-bit x86 architecture overview

- 8 general purpose registers `eax`, `ebx`, `ecx`, `edx`, `esi`, `edi`, `ebp`, `esp`
 - `esp` is the stack pointer
 - `ebp` is the frame pointer (optional)
 - Others are used for integer and pointer operations
 - 16- and 8-bit parts of the registers can be named (`ax` is least significant 16 bits of `eax`, `al` is least sig. 8 bits of `eax`, etc.)
- Instruction pointer `eip` holds the address of the next instruction to execute
- `eflags` register has bits like the zero flag or the carry flag that are set by arithmetic and logical operations, used for conditional control flow

Some x86 instructions (AT&T notation)

- `mov src, dest` ; Copies src to dest
- Arithmetic and bit operations
 - `add src, dest` ; computes $dest + src$, stores in dest
 - `sub src, dest` ; computes $dest - src$, stores in dest
 - `or`, `and`, `xor` all work the same way; `mul/div` use specific registers
- Stack operations
 - `push src` ; decrements `esp` by 4, writes `src` to stack
 - `pop dest` ; reads top of stack into `dest`, increments `esp` by 4

Some x86 instructions (AT&T notation)

- Function calls
 - `call foo ;` calls the function `foo`, pushes the address of the next instruction onto the stack
 - `leave ;` equivalent to `movl $ebp, $esp` followed by `popl $ebp`
 - `ret ;` pops the top of the stack into `eip` (returns from a function)
- Control flow
 - `cmp src2, src1 ;` computes `src1 - src2` and sets `eflags` register
 - `test src2, src1 ;` computes `src1 & src2` (bitwise-and) and sets `eflags`
 - `jz label ;` jump to label if the zero flag is set
 - `jnz label ;` jump to label if the zero flag is not set
 - `jc label ;` jump to label if the carry flag is set
 - `jnc label ;` jump to label if the carry flag is not set
 - `jmp label ;` unconditionally jump to label

Instruction suffixes

- l — (long) 32 bits
- w — (word) 16 bits
- b — (byte) 8 bits

- Examples
 - `movw %ax, %dx` ; Copies least sig. 16 bits of eax to least sig. 16 bits of edx
 - `pushl %edi`
 - `subl $16, %esp` ; Decrements esp by 16
 - `cmpl %edx, %eax` ; computes `eax - edx` and sets eflags based on the result

x86 operands

- Constants are prefixed with \$
- Registers are prefixed with %
 - `movb $8, %bl`
- Read/writing to memory has several forms
 - `(%eax)` ; Refers to the 1, 2, or 4 bytes at address stored in `eax`
 - `-8(%esp)` ; Address is `%esp - 8`
 - `4(%esi, %eax)` ; Address is `esi + eax + 4`
 - `16(%eax, %edx, 4)` ; Address is `eax + 4*edx + 16`

Using memory operands

- Load 4 bytes from `ebp + 4` into `eax`
 - `movl 4(%ebp), %eax`
- Store 1 byte from `dl` (least sig. 8-bits of `edx`) to address `edi`
 - `movb %dl, (%edi)`
- Add 4 bytes from address `edx` to `eax` and store in `eax`
 - `addl (%edx), %eax`
- Xor the constant `0x5555AAAA` with 4 bytes at address `8+ebp`
 - `xorl $0x5555AAAA, 8(%ebp)`

What values do `eax` and `edx` hold after this?

```
movl    $30, %eax
movl    $10, %edx
subl    %eax, %edx
addl    %eax, %eax
```

- A. `eax = 40, edx = 10`
- B. `eax = 60, edx = 40`
- C. `eax = 60, edx = -20`
- D. `eax = -40, edx = 10`

Function calls on 32-bit x86

- Stack grows down (from high to low addresses)
- Stack consists of 4-byte slots
- esp points to the bottom most “in-use” slot
- ebp “frame pointer” points to the previous ebp on the stack (if used)
- call pushes the return address onto the stack
- Function call arguments can be accessed at a positive offset from ebp
8(%ebp), 12(%ebp), 16(%ebp), etc.
- Local variables can be accessed at a negative offset from ebp
-4(%ebp), -8(%ebp), -12(%ebp), etc.

Warning!

- For most of these slides, the stack is drawn with low addresses on the bottom and high addresses on the top. The stack grows down both numerically and pictorially.

Function call example

```
1 int foo(int a, char *p) {  
2     int b = atoi(p);  
3     return a + b;  
4 }
```

eip →

```
1 foo:  
2     pushl    %ebp  
3     movl    %esp, %ebp  
4     subl    $40, %esp  
5     movl    12(%ebp), %eax  
6     movl    %eax, (%esp)  
7     call   atoi  
8     movl    %eax, -12(%ebp)  
9     movl    -12(%ebp), %eax  
10    movl    8(%ebp), %edx  
11    addl    %edx, %eax  
12    leave  
13    ret
```

esp →



← ebp

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esp →



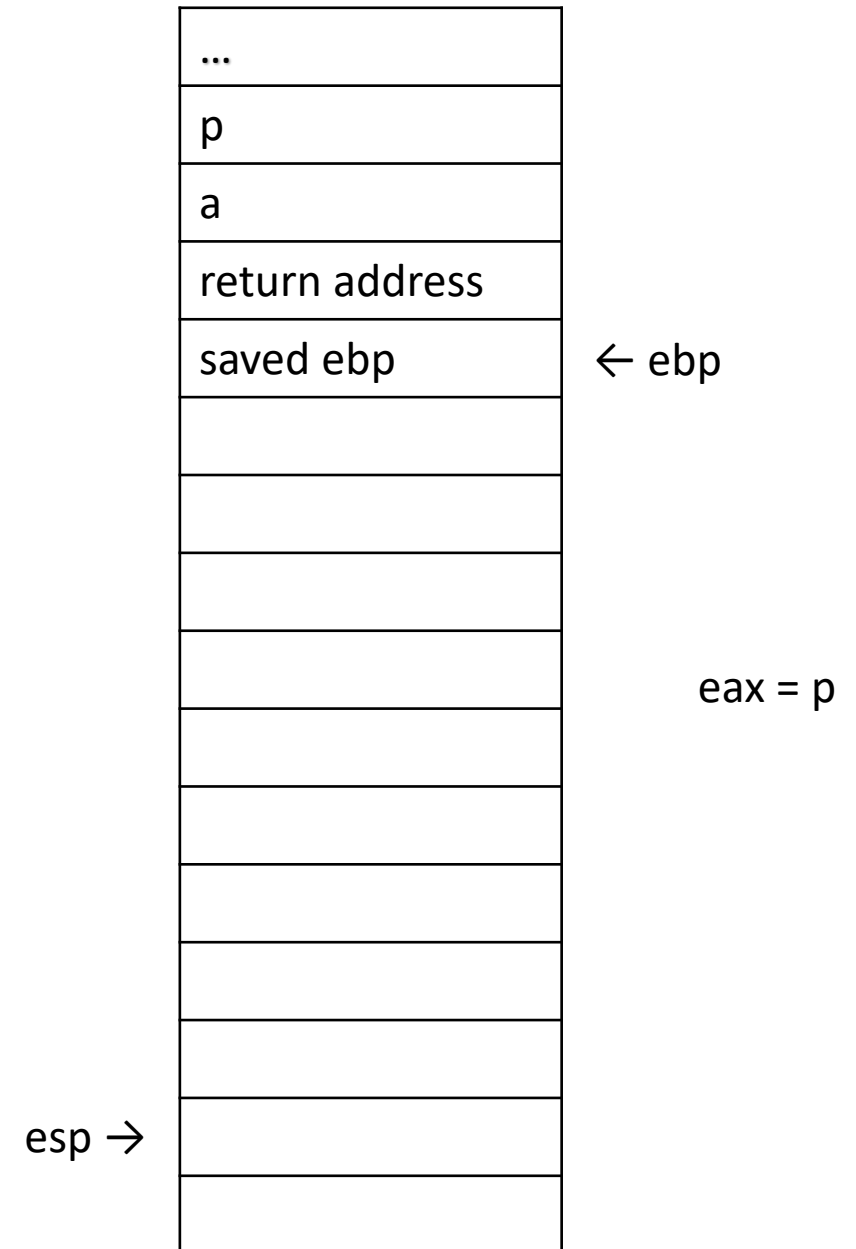
← ebp

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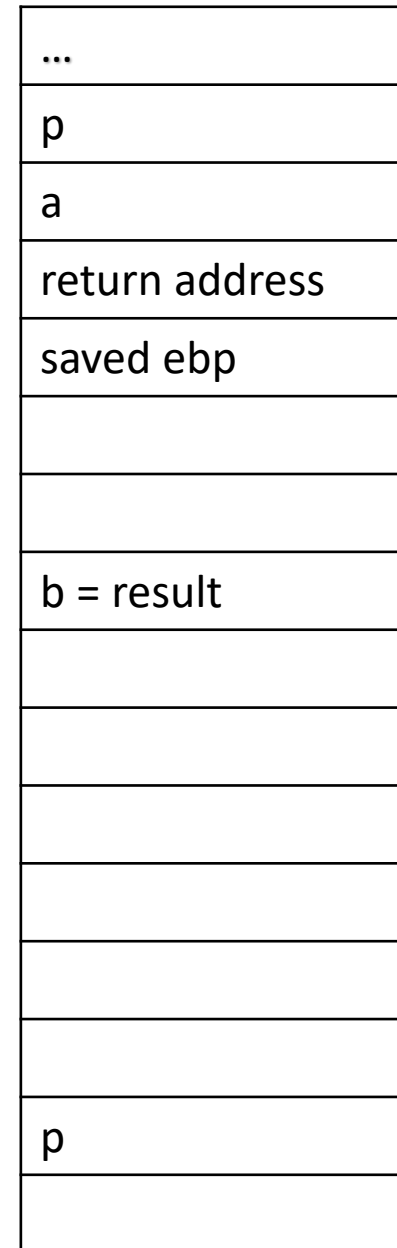


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```

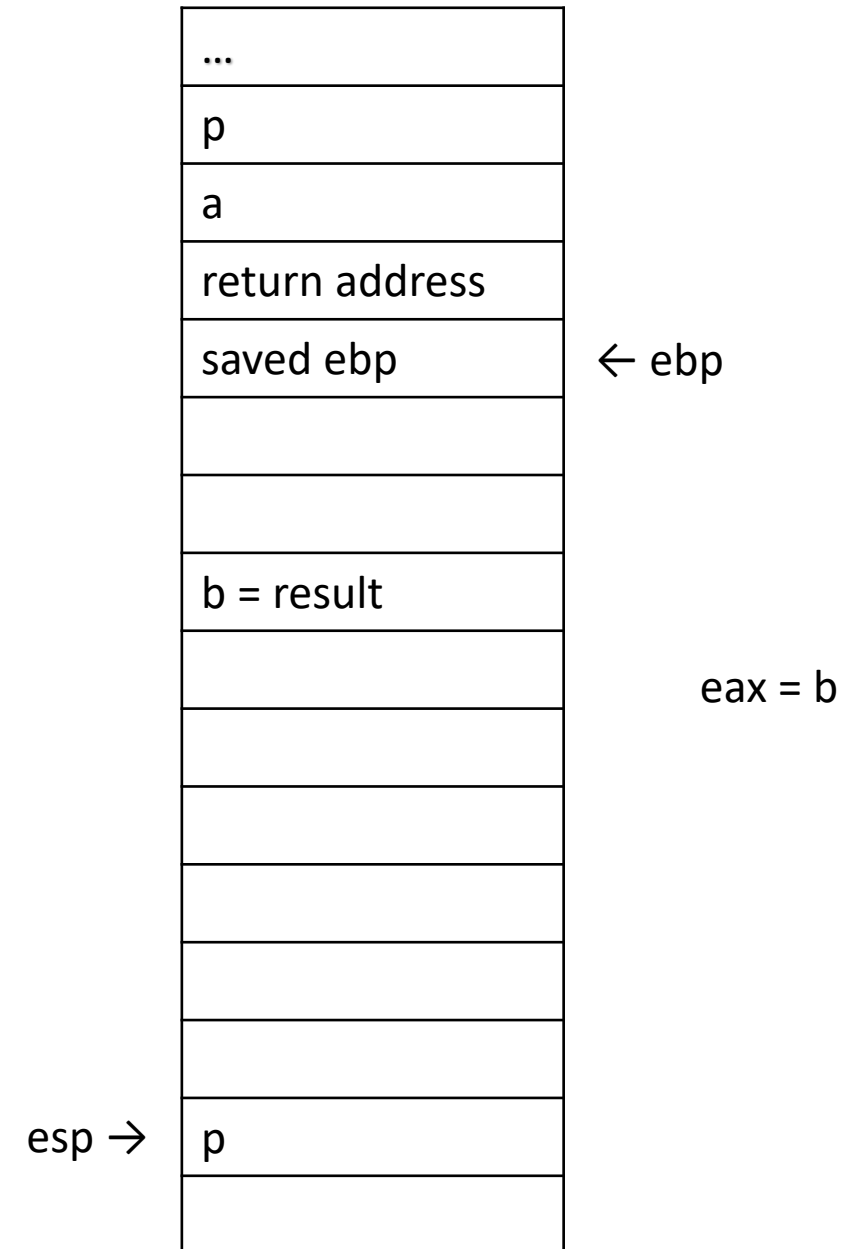
eip →



Function call example

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```

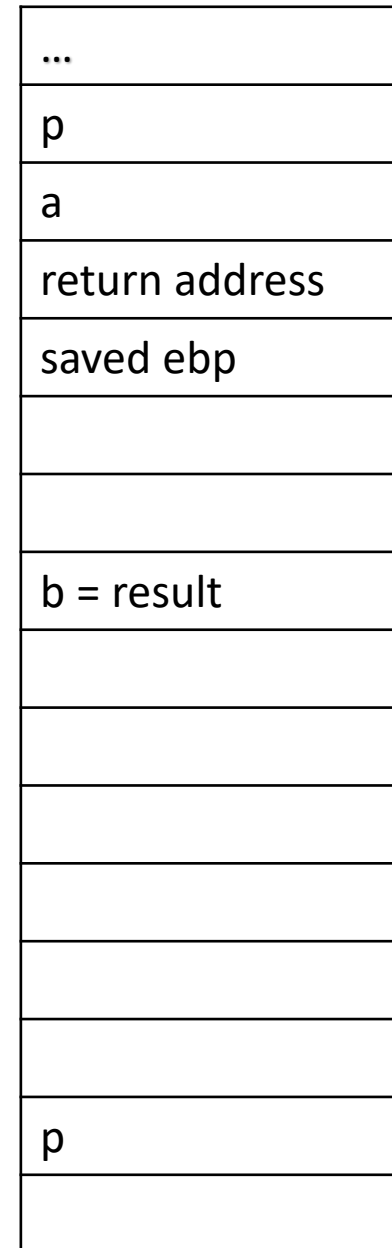


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12    leave  
13    ret
```

eip →



← ebp

eax = b
edx = a

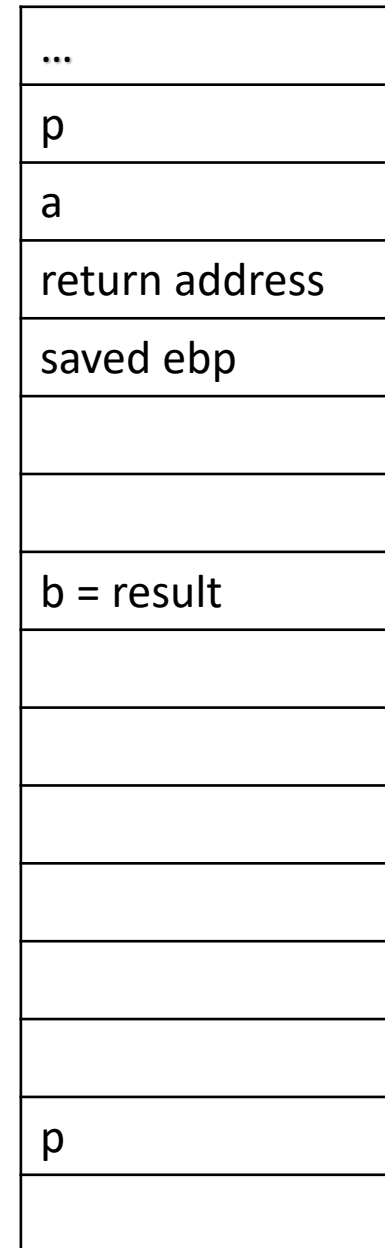
esp →

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12    leave
13    ret
```

eip →



eax = b + a
edx = a

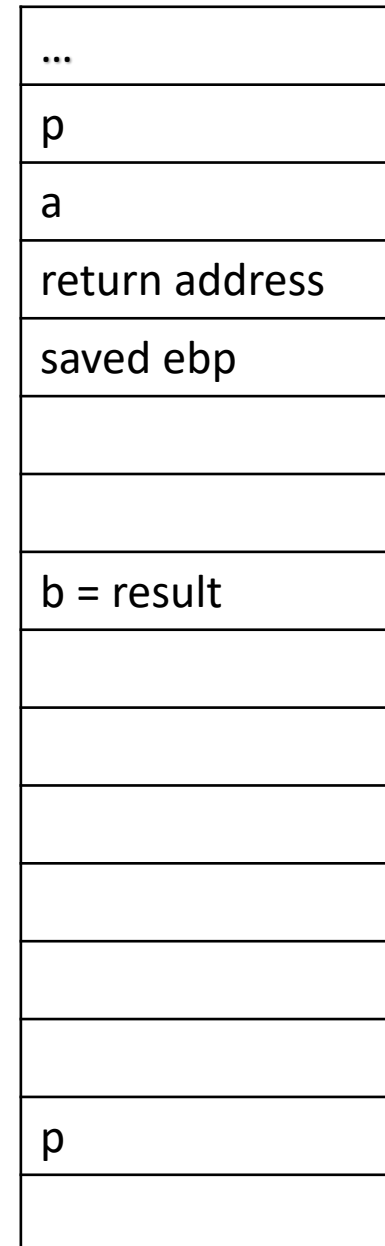
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12    leave  
13    ret
```

eip →

esp →



← ebp

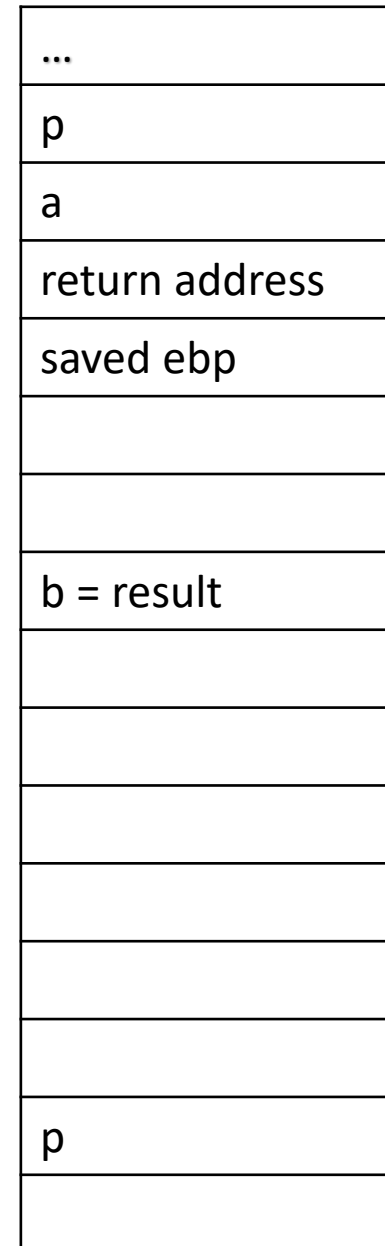
eax = b + a
edx = a

Function call example

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10    movl    8(%ebp), %edx
11    addl    %edx, %eax
12    leave
13    ret
```

esp →



← ebp

eax = b + a
edx = a
eip = ret addr