

# Homework 5

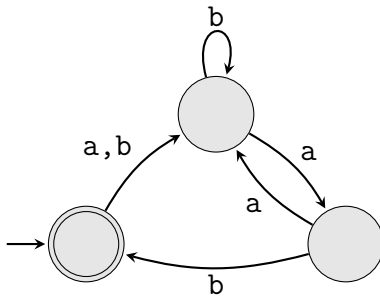
Due: Sunday, April 22, 2018

## Instructions

This assignment is due **Sunday, April 22, 2018 at 11:59PM (Central Time)**. Solutions must be submitted on Gradescope. Your solutions must be typeset. Handwritten solutions will not be graded and will receive a 0.

Late submissions will be accepted within **24 hours** after the deadline with a penalty of **25%** of the assignment grade. No late submissions will be accepted more than **24 hours** after the deadline.

**Problem 1** Answer the following questions about the following DFA  $M$  and give reasons for your answers.



- [2 points] Is  $\langle M, \text{abaab} \rangle \in A_{\text{DFA}}$ ?
- [2 points] Is  $\langle M, \text{bab} \rangle \in A_{\text{DFA}}$ ?
- [2 points] Is  $\langle M \rangle \in A_{\text{DFA}}$ ?
- [2 points] Is  $\langle M \rangle \in E_{\text{DFA}}$ ?
- [2 points] Is  $\langle M, M \rangle \in EQ_{\text{DFA}}$ ?

**Problem 2** Closure properties of decidable languages.

- [5 points] Prove that the class of decidable languages is closed under concatenation. [Hint: Let  $M_1$  and  $M_2$  be TMs that decide languages  $A$  and  $B$ . Construct a new TM  $M$  to decide  $A \circ B$ . TM  $M$  will take as input some string  $w = w_1w_2 \cdots w_n$  where each  $w_i \in \Sigma$  and will have to divide  $w$  into two pieces  $x = w_1w_2 \cdots w_k$  and  $y = w_{k+1}w_{k+2} \cdots w_n$  for some  $0 \leq k \leq n$  and check that  $x \in A$  and  $y \in B$ . Make sure that  $M$  tries all  $n + 1$  possible  $x$  and  $y$ .]

- b. [10 points] Prove that the class of decidable languages is closed under Kleene star. [Hint: Let  $M_1$  decide language  $A$ . Construct a new TM  $M$  to decide  $A^*$ . Recall that for string  $w$  to be in  $A^*$ , there must be a division of  $w$  into  $k$  pieces  $w = w_1w_2 \cdots w_k$  for some  $k \geq 0$  such that each  $w_i \in A$ . TM  $M$  will have to try all possible divisions for all values of  $k$  up to some number. If  $|w| = n$ , think about which values of  $k$  the TM  $M$  needs to consider. This problem is tricky!]

**Problem 3** Closure properties of Turing-recognizable languages.

- a. [10 points] Prove that the class of Turing-recognizable languages is closed under concatenation. [Hint: This is similar to the previous problem but now you have the issue that the TMs  $M_1$  or  $M_2$  may not halt on some division of  $w$  into  $x$  and  $y$ , but will halt and accept on some other division. Have  $M$  first write down all of the possible splits and then simulate  $M_1$  and  $M_2$  on each of the possible  $x$  and  $y$  in “parallel” by performing one step of the simulation of each TM at a time.]
- b. [15 points] Prove that the class of Turing-recognizable languages is closed under Kleene star. [Hint: Use the hints for Problems 2b and 3a.]

**Problem 4** [10 points] Let  $A = \{\langle R, S \rangle \mid R \text{ and } S \text{ are regular expressions such that } L(R) \subseteq L(S)\}$ . Prove that  $A$  is decidable by giving a TM that decides it. [Hint: Your TM should construct one or more DFAs and use a decider for a language we’ve already shown to be decidable as a subroutine.]

**Problem 5** [10 points] Let  $B = \{\langle R \rangle \mid R \text{ is a regular expression describing a language containing at least one string } w \text{ that has } 111 \text{ as a substring (i.e., } w = x111y \text{ for some } x \text{ and } y)\}$ . Prove that  $B$  is decidable by giving a decider for it. [Hint: Your TM should construct one or more DFAs and use a decider for a language we’ve already shown to be decidable as a subroutine.]

**Problem 6** [10 points] Prove that  $EQ_{CFG}$  is undecidable. [Hint: Give a reduction from  $ALL_{CFG}$ .]

**Problem 7** [10 points] Let  $C = \{\langle M \rangle \mid M \text{ is a TM that accepts } w^R \text{ iff it accepts } w\}$ . Prove that  $C$  is undecidable. [Hint: Give a reduction from  $A_{TM}$ .]

**Problem 8** [10 points] Consider the problem of determining whether a Turing machine  $M$  on an input  $w$  ever attempts to move its head left when its head is on the left-most tape cell. Formulate this problem as a language and show that it is undecidable.