

CSCI 275: Programming Abstractions

**Lecture 11: Higher Order Wrap-Up
Spring 2025**

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Slides gratefully borrowed from Molly Q Feldman**



Questions? Comments?

Functional Language of the Week: Python

- Wait, hold on! Python is *not* a functional paradigm language
 - Paradigms are a gray area.....
- The transition from Python 2 to Python 3 facilitated significantly better functional programming in Python
 - This was, more or less, due to popular demand
- A (long-running) guide to functional programming in Python
<https://docs.python.org/3/howto/functional.html#>



Functional Language of the Week: Python

```
#filter
```

```
x = filter(lambda x: x%3 == 0 and x%5 != 0, [3,6,9,12,15,18,21,24,27,30])  
print(list(x)) #gives [3, 6, 9, 12, 18, 21, 24, 27]
```

```
#map
```

```
a = map(lambda x: x + 1, [1,2,3])  
print(list(a)) #gives [2, 3, 4]
```

```
#subtle: how do you change how sorted works?
```

```
#a lambda!
```

```
r = sorted([(1,0), (2,1), (3,2)], key = lambda x: x[1])  
print(list(r)) #gives [(1, 0), (2, 1), (3, 2)]
```



foldl

Reminder: Light Switch State Machine

Possible actions: 'up, 'down, 'flip

Possible states: 'on, 'off

We want

(state-after ' (up flip)) => 'off

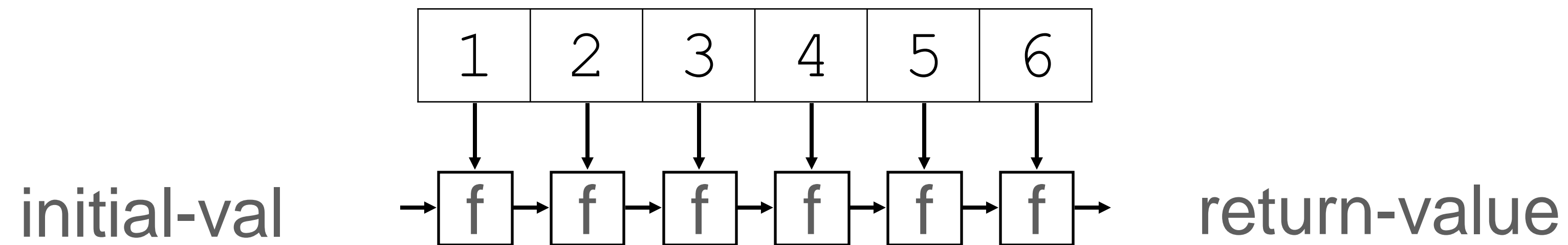
Reminder: Lightswitch, as `foldr` and `foldl`

```
(define state-after  
  (lambda (actions)  
    (foldr next-state 'off (reverse actions))))
```

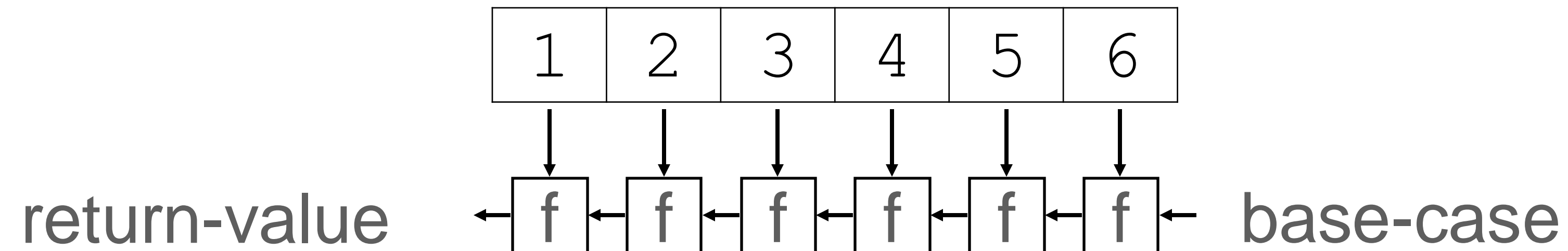
```
(define (state-after-left actions)  
  (foldl next-state 'off actions))
```

foldl vs. foldr

`foldl` combines elements of the list starting with the first (left-most) element



`foldr` combines elements of the list starting with the last (right-most) element



But wait: more “thoughtful” motivation for foldl

Reminder: Tail Recursion and using an “accumulator”

```
(define (fact-a n acc)
  (if (<= n 1)
      acc ; return the accumulator
      (fact-a (sub1 n) (* n acc))))

(define (fact2 n)
  (fact-a n 1))
```

Four things to notice:

- We defined a recursive helper function that takes an **additional param**
- We provide an **initial value** for the accumulator in `fact2`'s call to `fact-a`
- The base case returns the **accumulator**
- `fact-a` is tail-recursive

Product: An Accumulator Pattern

```
(define (product-a lst acc)
  (cond [(empty? lst) acc]
        [else (product-a (rest lst)
                          (* (first lst) acc))]))
```

```
(define (product lst)
  (product-a lst 1))
```

Reverse: An Accumulator Pattern

```
(define (reverse-a lst acc)
  (cond [(empty? lst) acc]
        [else (reverse-a (rest lst)
                          (cons (first lst) acc))]))

(define (reverse lst)
  (reverse-a lst empty))
```

Map: An Accumulator Pattern

```
(define (map-a proc lst acc)
  (cond [(empty? lst) acc]
        [else (map-a proc (rest lst)
                      (cons (proc (first lst)) acc))]))

(define (map proc lst)
  (reverse (map-a proc lst empty)))
```

Accumulator Pattern Similarities

Basic structure is the same (rewriting slightly)

```
(define (fun-a lst acc)
  (cond [(empty? lst) acc]
        [else
         (fun-a (rest lst)
                 (combine (first lst) acc))]))
```

```
(define (fun ... lst)
  (fun-a lst initial-val))
```

Function	initial-val	(combine head acc)
product	1	(* head acc)
reverse	empty	(cons head acc)
map	empty	(cons (proc head) acc)

We must reverse the result

Abstraction: fold left

`(foldl combine initial-val lst)`

`combine: $\alpha \times \beta \rightarrow \beta$`

`initial-val: β`

`lst: list of α`

`foldl: $(\alpha \times \beta \rightarrow \beta) \times \beta \times (\text{list of } \alpha) \rightarrow \beta$`

Elements of `lst` = $(x_1 \ x_2 \ \dots \ x_n)$ and `initial-val` are combined by computing

$z_1 = (\text{combine } x_1 \ \text{initial-val})$

$z_2 = (\text{combine } x_2 \ z_1)$

$z_3 = (\text{combine } x_3 \ z_2)$

\vdots

$z_n = (\text{combine } x_n \ z_{n-1})$

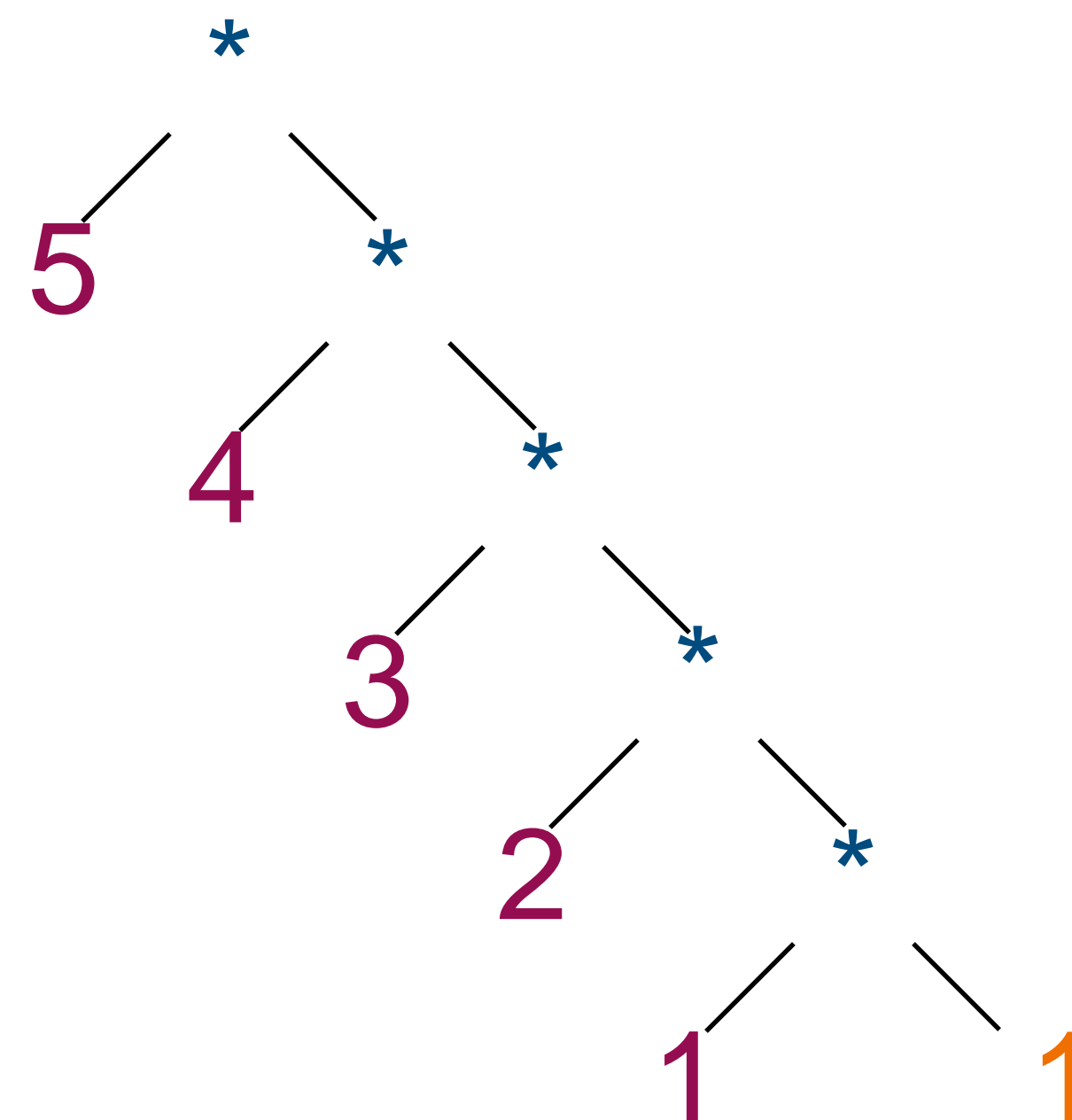
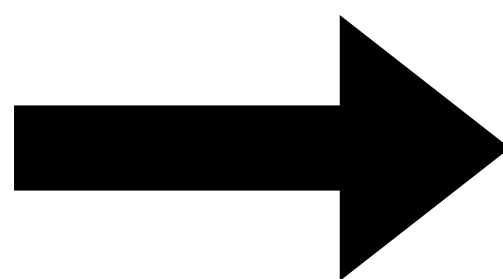
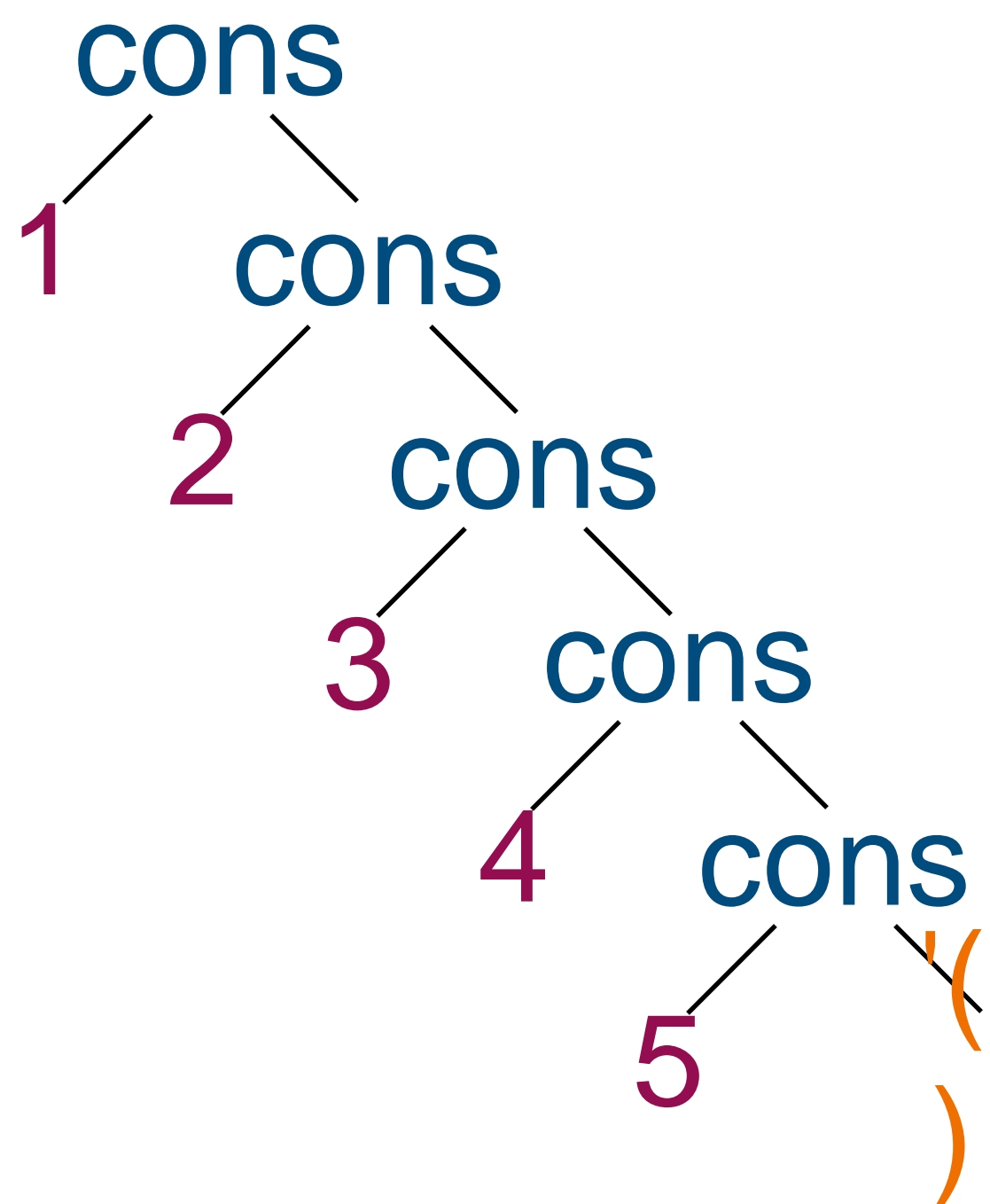


product as fold left

`(foldl combine initial-val lst)`

```
(define (product lst)
  (foldl * 1 lst))
```

`combine: number × number → number`
`initial-val: number`
`lst: list of number`

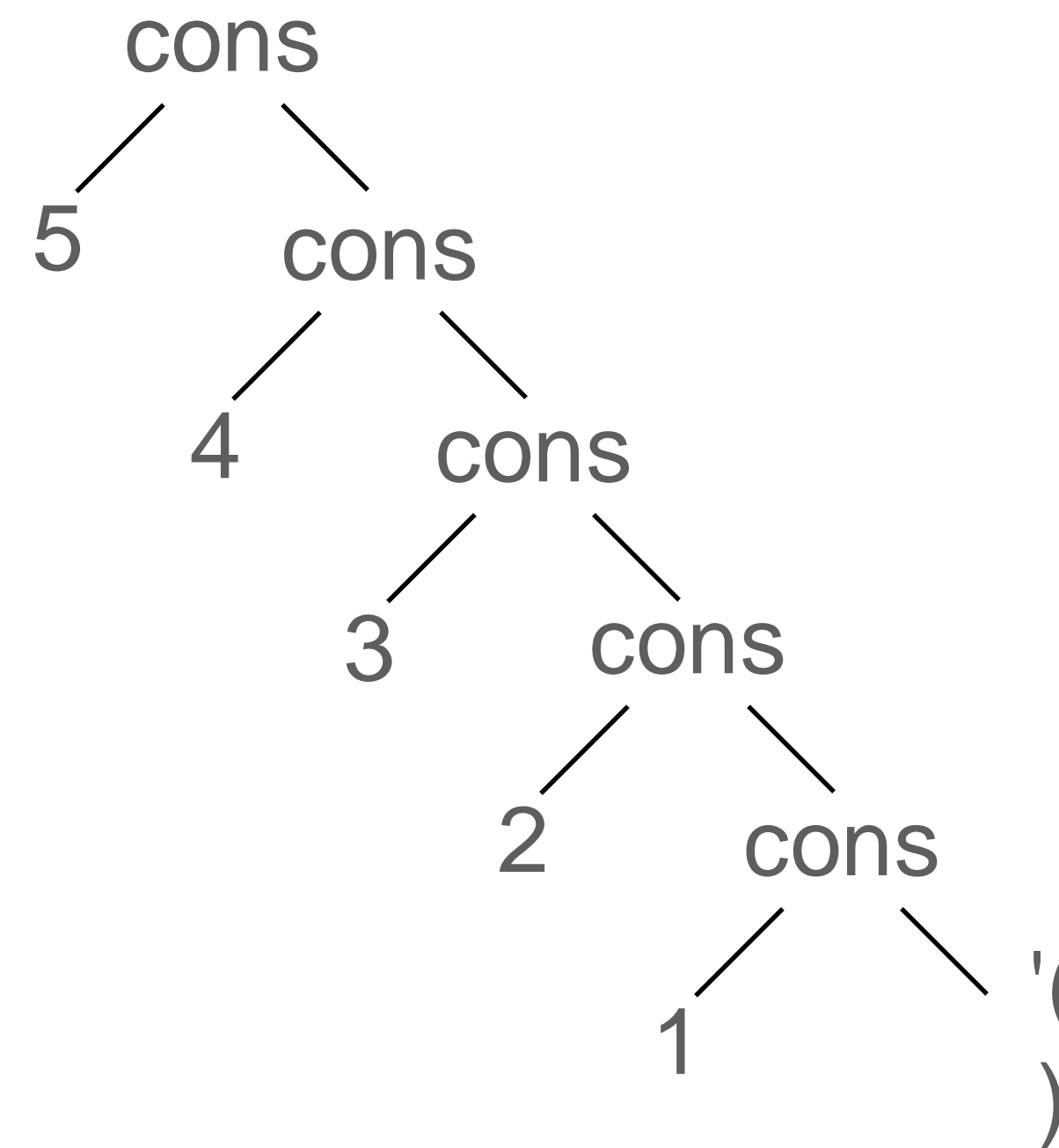
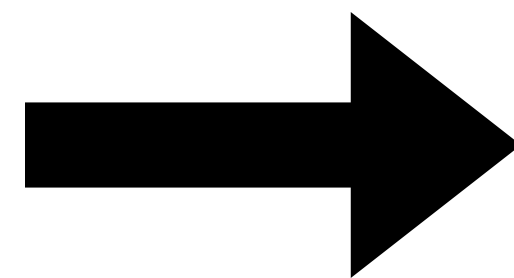
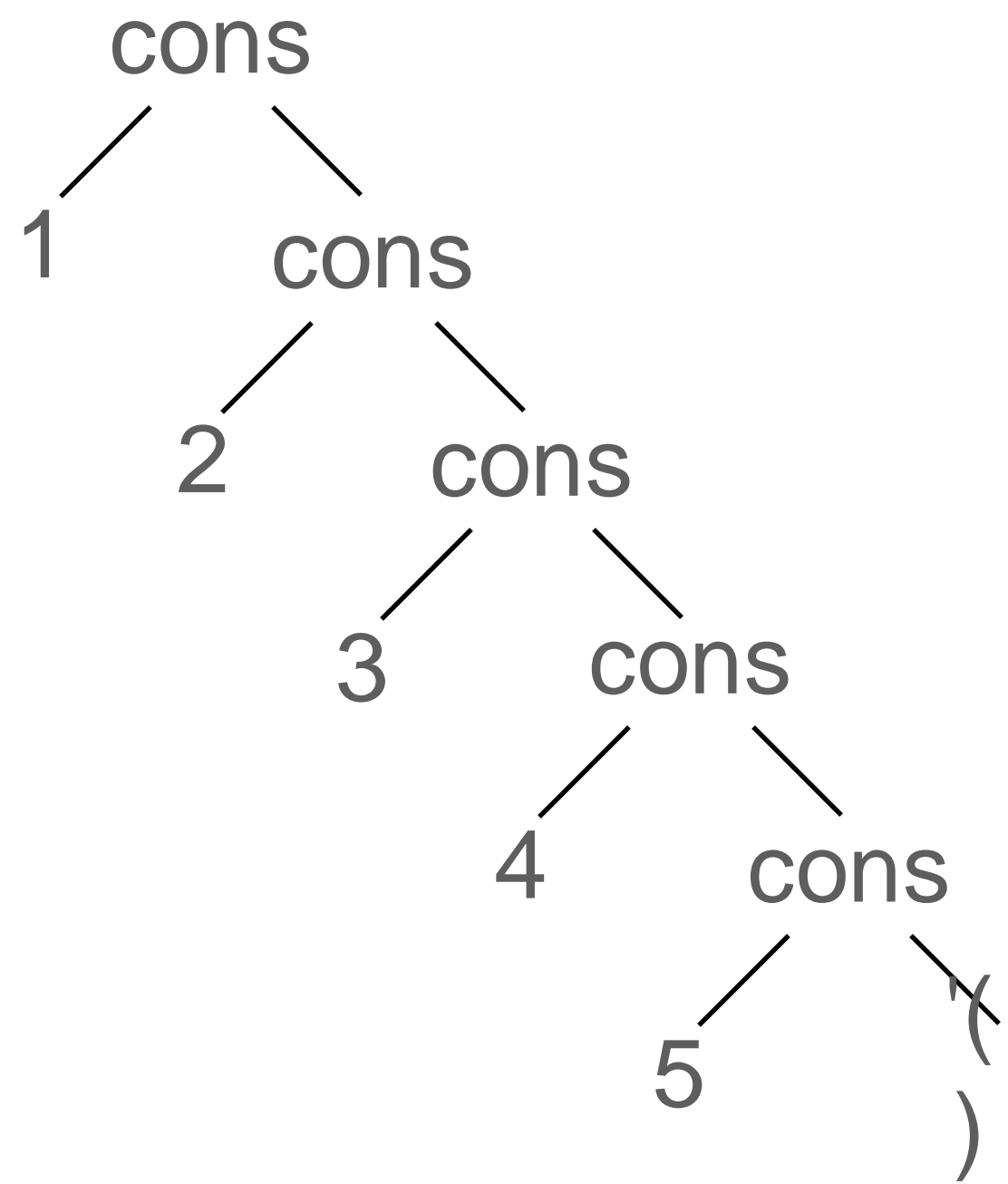


reverse as fold left

(foldl combine base-case lst)

```
(define (reverse lst)
  (foldl cons empty lst))
```

combine: $\alpha \times \text{list of } \alpha \rightarrow \text{list of } \alpha$
initial-val: list of α
lst: list of α



Which fold to pick?

- “Most of the time”, either will work +/- a call to `reverse`
- Be careful when `combine` has ordering effects
- If the computation makes more sense as a right-to-left computation on the elements of the list, then use `foldr`
- But, most of the time, use `foldl`
 - Lists run left-to-right in Racket world
 - Fold in most other functional contexts assumes `foldl`
 - Tail recursive, and thus more efficient

Aside: foldr can work on lazy, infinite lists

Languages with lazy evaluation can have infinite data structures like infinite-length lists (we'll see something similar in Racket later in the semester)

Fold right can work with infinite lists so long as only a finite portion of the list is required to compute the value

```
nats :: [Int]
nats = [1..] -- An infinite list
```

```
sumFirst :: Int -> Int
sumFirst n = foldr (\x acc -> if x <= n
                              then x + acc
                              else 0)
                (error "this will never be evaluated")
                nats
```

Variable Argument Procedures

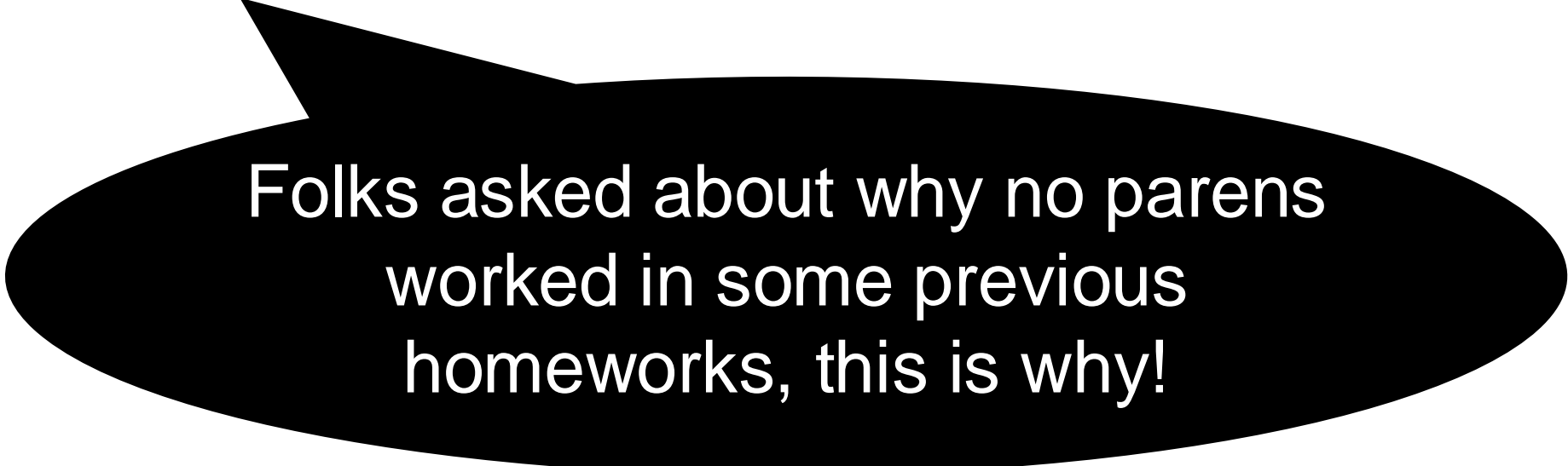
Variable argument procedures

```
(define foo (lambda (params) body))
```

When `params` is a **list of identifiers** (as we know it thus far!), the identifiers are bound to the values of the procedure's arguments

When `params` is an **identifier** (i.e., not a list), then the identifier is bound to a list of the procedure's arguments

```
(define count-args  
  (lambda (params)  
    (length params)))  
  
(count-args 'a 2 #f) => 3
```



Folks asked about why no parens worked in some previous homeworks, this is why!

Required parameters + variable parameters

```
(define foo (lambda (x y z . params) body))
```

Separate the required parameters from the list of variable parameters with a period

```
(define drop-2  
  (lambda (x y . lst) lst))
```

```
(drop-2 1 2 3 4)
```

x is bound to 1

y is bound to 2

lst is bound to '(3 4)

Review & Practice

There's a standard library procedure (`round x`) that takes a number as input and rounds it to the nearest integer.

If we have a list of numbers `'(1.1 2.9 3.5 4.0)` and we want a list of rounded numbers `'(1.0 3.0 4.0 4.0)`, how can we get that?

A. `(map (round x) '(1.1 2.9 3.5 4.0))`

B. `(map (lambda (x) (round x)) '(1.1 2.9 3.5 4.0))`

C. `(map round '(1.1 2.9 3.5 4.0))`

D. `(round '(1.1 2.9 3.5 4.0))`

E. More than one of the above

Distance of a 2-d point from the origin

Recall that a point (x, y) lies $\sqrt{x^2 + y^2}$ from the origin

Let's make a procedure to compute this

```
(define (distance-from-origin x y)
  (sqrt (+ (* x x) (* y y))))
```

```
(distance-from-origin 3 4) => 5
```

```
(define (distance-from-origin x y)
  (sqrt (+ (* x x) (* y y))))
```

If we have a point

(define p '(5 -8)) how can we get its distance from the origin?

A. (distance-from-origin p)

B. (apply distance-from-origin p)

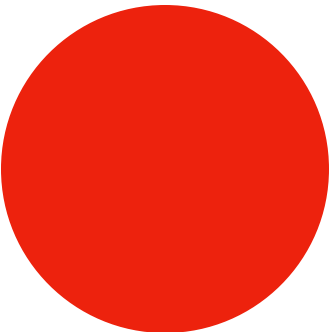
C. (distance-from-origin (first p) (second p))

D. More than one of the above

Shapes

Racket library `2htdp/image` has procedures for creating images

```
(require 2htdp/image)
```

```
(circle 20 'solid 'red) => 
```

```
(rectangle 50 20 'outline 'blue) => 
```

radius

width **height**

If we have a list of radii, say `lst` is `(20 30 50 60)` and we want a list of solid, red circles with those radii, which should we use?

`(_____ lst) => (list )`

A. `(map circle 'solid 'red lst)`

B. `(map (lambda (r) (circle r 'solid 'red)) lst)`

C. `(apply circle 'solid 'red lst)`

D. `(apply (lambda (r) (circle r 'solid 'red)) lst)`

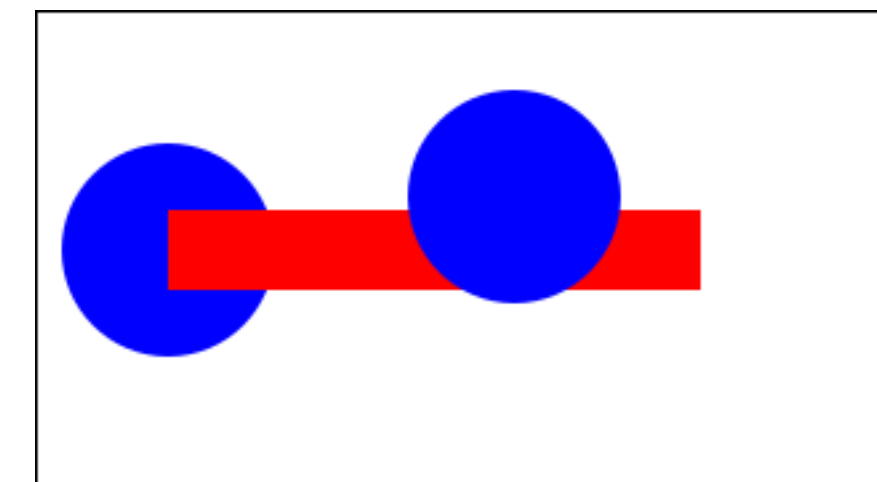
E. `(foldr (lambda (r) (circle r 'solid 'red)) empty lst)`

Combining images

`(empty-scene 320 180)` gives a white rectangle with a black border we can draw on

`(place-image img x y scene)` returns a new image by starting with `scene` and drawing `img` at `(x, y)`

```
(let* ([c (circle 40 'solid 'blue)]
       [r (rectangle 200 30 'solid 'red)]
       [s0 (empty-scene 320 180)]
       [s1 (place-image c 50 90 s0)]
       [s2 (place-image r 150 90 s1)]
       [s3 (place-image c 180 70 s2)])
  s3)
```



Imagine we have a list of 3-element lists (shape x y), e.g., lst is the list

```
(list (list (circle 40 'solid 'blue) 50 90)
      (list (rectangle 200 30 'solid 'red) 150 90)
      (list (circle 40 'solid 'purple) 180 70))
```

How would you draw those shapes on a scene at their coordinates?

- A.

```
(map (lambda (i)
      (place-image (first i) (second i) (third i) scene))
     lst)
```
- B.

```
(apply (lambda (i)
      (place-image (first i) (second i) (third i) scene))
       lst)
```
- C.

```
(foldr (lambda (i s)
      (place-image (first i) (second i) (third i) s))
       scene
       lst)
```

**Try out the previous question
on your own!**