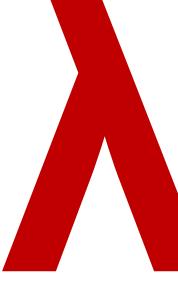
CSCI 275: **Programming Abstractions** Lecture 02: Procedures & Choice Spring 2025

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Announcements

Homework 0 is up on the website

Goals for Today

- Basics of Racket
- How do we make choice (i.e., conditionals, etc.)?
- How do we construct and use procedures?

Introducing Racket

When we talk about code/Racket in this class, I will do my best to use Font in This Text to differentiate what is description and what is code PowerPoint will mess up my "quotation marks"





Why Racket for CS 275?

All LISP-type languages have lists as the main data structure

- Programs are lists
- Data are lists
- Racket programs can reason about other programs. This makes Racket useful for thinking about programming languages in general.

Racket is a different programming paradigm

- Python, Java, C and other languages are imperative languages. Programs in these languages do their work by changing data stored in variables Racket programs can be written as functional programs—they compute by evaluating functions and avoid variable assignments.

Why Racket for CS 275?

Racket is very elegant. It is much less verbose than Java, for instance, which means it is easier to see what is happening in a Racket program.

I think its fun.

It lets you learn functional programming without a lot of extra features.

Racket Basics

We are used to **basic values** in most languages

- Numbers (Integers & Floats)
- Strings
- Booleans

These also can look different depending on the language! 'banana' is invalid Java, but valid Python

We are also accustomed to procedures/functions which act on elements of these types





Arithmetic/logical/string operations

3 + 5 x • (4 + y + z) x AND y x OR y OR z "hello" + " " + "world" (

Language Design Statement: you know the *semantics* of these terms, even if this *syntax* is not that of a language you've learned before

Everything is prefix in Racket

Language Design Statement:

arbitrary.

In Racket, you put the operator or function call *first* (prefix form)

(< x 2) instead of x < 2

The order that a language has the operators and operands is

Equivalent operations in Racket

- 3+5 (+ 3 5)
- $x \cdot (4 + y + z)$ (* x (+ 4 y z))
- x AND y (and x y)
- x OR y OR z (or x y z)

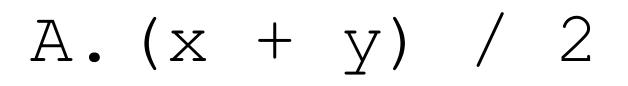
"hello" + " " + "world" (string-append "hello" " " "world")

Some basic data types and variables in Racket

- Numbers: 83, -6, 25.23
- Strings: "this is a string"
- Booleans: #t, #f
- We can define variables* using define: (define department "CSCI") (define course-number 275) (define instructor "Stephen Checkoway")

* These are not really variables; they're just names we give to some values

In most languages, we would compute the arithmetic mean (average) of two numbers (or variables holding numbers) as (x + y) / 2. How do we do this in Racket?



- B.((x + y) / 2)
- C. (+ x y / 2)
- D. (+ (/ x y) 2) E. (/ (+ x y) 2)

What do you think these examples will evaluate to?

(+ 5 2) (zero? x) (or (and #t #f) (and #t #t)) (+ (- 1 0) (- 2 3))

What do you think these examples will evaluate to?

(+ 5 2) (zero? x)(or (and #t #f) (and #t #t)) #t (+ (- 1 0) (- 2 3))

Depends on x

Procedures in Racket

All the examples we saw on the previous example - e.g. (zero? X) and (+ (- 1 0) (- 2 3)) - are calls to procedures

In general, the structure of a procedure call in Racket is: <mark>(</mark>name-of-procedure argl arg2 … argn<mark>)</mark>

The parentheses here are the call to name-of-procedure The arguments are given after the procedure's name, separated by spaces

Procedure calls and special forms

first element of the sequence (here, foo)

etc.)

the arguments and returns the result

Otherwise, error!

(1 2 3) is an error because 1 is not special form or procedure

- When presented with a sequence (foo arg1 arg2 ...) Racket looks at the
- If foo is a special form, Racket follows special instructions (define, and,

If foo is a procedure (built-in or made by you), it applies that procedure to

This is the most common error in the first couple weeks of class!





Special Form: define

(define id s-exp) The define special form binds an identifier to a value

(define hi "Hello") (third professors) => "Cynthia"

(define x (+ 20 100))



Giving names to values – useful! However, these are not variables.

- This modifies the *environment*, the mapping of identifiers to values

(define professors '("Molly" "Steve" "Cynthia"))

Whatever is in s-exp is evaluated, so x is bound to 120









Predicates



Racket has a bunch of procedures that return #t if its argument satisfies some property

(zero? x) (empty? x) (positive? X) (number? x)

Style: predicates in Racket will always have ? as the last character (they are asking a question!)

returns #t if x is equal to 0 returns #t if x is the empty list returns #t if x is a positive number returns #t if x is a number



Tests for equality

Most of the time: Use equal?

(equal? a b) compares structures recursively

Are you dealing with numbers? Use = (= a b) compares only numbers, cannot be used for anything else

eq?/eqv? are about referring to the same object in memory; sometimes useful when you care about literal equality



If expression

(if test-exp then-exp else-exp)

expression evaluates to the evaluation of then-exp

If test-exp evaluates to #f, then the whole if expression evaluates to the evaluation of else-exp

(if (= x y))(+ x 2) V)(if (empty? lst) "The list is empty" "The list is not empty")

If test-exp evaluates to anything other than #f, then the whole



Conditional expressions (cond [test-exp1 exp1] ... [test-expn expn])

Evaluates the test-exp expressions in turn

whole expression

We can (and should!) use else as the last test expression

(cond [(zero? x) 0] [(> x 0) 1] [else -1])

The first one that evaluates to something other than # f has its corresponding exp evaluated - this becomes the value of the

If your program is more than just a *very simple* if statement, use cond. It's good style.



(define foo 12) (cond [(< foo 2) #t] [(>= foo 10) #f] [(not (zero? foo)) #t]

What does this code evaluate to?

- A.#t
- B.#f
- C.#t or #f, depends on the run
- D.Error
- E.Something else

[else (error "there is a problem!")])



Some questions (cond [(< foo 2) #t] [(>= foo 10) #f]

- 1. How can I get the cond to take an argument, rather than just reference a "global" foo?
- 2. How do I "save" code like that above to be able to reuse it? (i.e. a function!)
 - How is/isn't this related to using define to bind identifiers?

(define foo 12) [(not (zero? foo)) #t] [else (error "there is a problem!")



Creating procedures: lambda Procedures are creating using the lambda special form

(lambda parameters body ...)

is called

procedure, they're evaluated in turn

parameters is an unevaluated list of identifiers which will be bound to the values of the procedure's arguments when procedure

body is a sequence of s-expressions that form the body of the (lambda (x y) (/ (+ x y) 2))(lambda (name) (displayln "Hello ") (displayln name))



Naming lambdas Given we have a lambda, we can use it and call it

- This will evaluate to 6. However, this current structure doesn't allow us to reuse the lambda with a different input.
- We already have a way to bind a value to an identifier ("name"): that's define.
- We know define attaches a name to an evaluated value (define x (+ 20 100)) means x is bound to 120

- ((lambda (x) (+ x 2)) 4)

So what does a lambda evaluate to? Anything?

BIG IMPORTANT SLIDE

Unlike procedures in most languages, in Racket there is a notion that lambdas are values & so can be evaluated

- lambdas are like numbers, strings, lists, etc.
- We can pass them around, return them, hold them as their own, evaluated concept
 - This is really not true in languages like C, for instance • This makes procedures first-class in Racket
- Support for higher-order/first-class functions is one of the hallmarks of a language that supports functional programming



Closures: what lambdas evaluate to

The expression of (lambda parameters body...) evaluates to a *closure* consisting of

- The parameter list (a list of identifiers)
- The body as un-evaluated expressions (often just one expression)
- time the lambda expression is evaluated

We'll return to this becomes important! - The environment (the mapping of identifiers to values) at the



define + lambda = reusable procedures! We can combine define and lambda, so that we can get a named procedure!

(define add-two (lambda (x) (+ x 2)))

To call it, we then use prefix call notation, as usual:

- (add-two 2) will give us 4

What have we learned thus far?

- How to call procedures
- Predicates
- if
- cond
- define
- lambda
- define & lambda together!

(define lily (lambda (x y) (string-append y x)))

(lily "hello" "?")

What does this code evaluate to?

- A.Error
- B. "hello?"
- C."?hello"
- D. "hello ?"
- E.Something else

(define alright (lambda (a b) (cond [(equal? a b) "equal"] [(positive? a) 17] [else "chaos!"])))

What does calling (alright 10 -30) evaluate to?

- A. "chaos"
- B.Error
- C.5
- D.17
- E."equal"

- [(and (positive? a) (negative? b)) 5]

Can we use identifiers in lambdas? Sure!

Note: you won't see for loops in this class; recursion all the way

Computing factorial in Racket: (define fact (lambda (num) (if (<= num 1) 1 (* num (fact</pre>

(* num (fact (- num 1))))))

A Note on Readings RPTFW is really a reference guide

- If something didn't make sense in lecture? Great resource, this textbook or the additional resources I link • Honor Code: look it up there, not Google!
- If you want more detail about something? Readings!
- Especially Chapters 1 & 2 teach you about some great Racket operators (hint: member, remove) that we don't cover in class
- You'll read about mutability (e.g., set!), for loops and some "useful" Racket that is not functional style - refrain from using it and stick to what we learn in class!
- Readings/order of lecture not entirely in sync

Next Up! See the Schedule for Readings. Homework 0 is live - If you've never used Git/Github locally, please start

- ASAP
- Due Friday at 23:59

Post on Ed with questions