CSCI 275: Programming Abstractions Lecture 20: MiniScheme D, Conditionals Fall 2024

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Some General Purpose Notes

Racket Style

- Never use multiple conds, when you can just use one!
- Nested defines? letrec instead.
- Logic choices should be *distinct* in a cond

- When building something recursive, base cases first in the cond - else case should almost always be recursive call, not #f



Testing

- Generally, *always* good to te that operates over a list
- MiniScheme testing is going to be intense!
 Lots of tests, best way to check that it works
- Learning how to write a good test suite: part of the class
 Think about "code coverage"
 Every branch of a cond should have a related test

- Generally, always good to test the empty list case for anything

Summary Problems

- Available!
- Your final is based strongly on these problems - In your interest to get feedback early part of your grade on the Final Project
- Style, how you solve the problems, and test the solutions are

Back to MiniScheme

What have we discussed already? (A-C)

 $EXP \rightarrow number$ symbol parse into var-exp (EXP EXP*) parse into app-exp

- Parsing and evaluating numbers Parsing and evaluating symbols • Parsing and application (primitive functions)

- What an environment does

parse into lit-exp

MiniScheme D: Conditionals

Heads Up! MiniScheme Booleans & If

Booleans in MiniScheme are different than in Scheme/Racket!

Booleans in MiniScheme are True and False False to init-env Bind them to 'True and 'False

We'll treat anything other than False and 0 as being true for conditionals.

- Like our primitive procedures, you'll need to add symbols True and

What value does MiniScheme return for this expression assuming that x is bound to 23 and y is bound to 42? (if (- y x) 25 37)

A. 25

B. 37

C. It's an error because (-y x) is a number

Our first special form: if

EXP → number parse into lit-exp symbol parse into var-exp (if EXP EXP EXP) parse into ite-exp (*EXP EXP*^{*}) parse into app-exp

We need a new data type for the if-then-else expression: ite-exp

- ite-exp
- ite-exp?
- ite-exp-cond
- ite-exp-then
- ite-exp-else

Parsing special forms if, let, lambda, etc.

(define (parse input) (cond [(number? input) (lit-exp input)] [(symbol? input) (var-exp input)] [(list? input) (cond [(empty? input) (error ...)] [(eq? (first input) 'if) ...] Make sure that [(eq? (first input) 'let) ...] input is the right length for each [(eq? (first input) 'lambda) ...] special form! • • • [else (app-exp ...)])] [else (error 'parse "Invalid syntax ~s" input)]))

Parsing if-then-else expressions

If-then-else expressions are recursive

- E.g., $EXP \rightarrow (if EXP EXP EXP)$
- When parsing an if-then-else expression, you want to parse the sub expressions using parse
- The input to parse will look like '(if (lt? x 1) (+ y 100) Z)
- The condition is (second input) The then-branch is (third input)
- The else-branch is (fourth input)

Evaluating ite-exp

- Parse tree is recursive: (parse '(if x 10 20)) (ite-exp (var-exp 'x) (lit-exp 10) (lit-exp 20))
- When evaluating, you should call eval-exp recursively
 - First, call it on the conditional expression
 - If the condition evaluates to False or 0,
 - evaluate the last expression and return its result
 - herwise.
- evaluate the middle expression and return its result

What happens if you implement eval-exp for an ite-exp by calling eval-exp on all three parts of the expression before deciding which one to return? (let ([co (eval-exp (ite-exp-cond tree) e)] [th (eval-exp (ite-exp-then tree) e)] [el (eval-exp (ite-exp-else tree) e)]) (if co th el))

- A. The code works perfectly

- D. The code will produce the wrong results on all inputs

B. The code works correctly, but inefficiently on some inputs (which?)

C. The code will produce the wrong result on some inputs (which?)