Programming Abstractions

Week 11-2: MiniScheme H

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What is the result of this expression?

```
(let ([f (λ (n)
           (if (= 0 n)
               empty
               (cons n (f (- n 1)))))])
  (f 4))
```

A. `'(0 1 2 3 4)`

B. `'(1 2 3 4)`

C. `'(4 3 2 1 0)`

D. `'(4 3 2 1)`

E. An error

Implementing recursion in MiniScheme H

```
(letrec ([f exp1] [g exp2] ...) body)
```

We'll have the parser parse a `letrec` expression into something equivalent that uses only things we have implemented

We won't need to change `eval-exp` at all!

Two options

We can use the Y combinator (technically the Z combinator)

We can use `set!/begin`

Which would you prefer?

Z combinator it is!

```
Z = λf.(λx.f(λ v.xxv))(λx.f(λv.xxv))
```

Translated from λ-calculus to Scheme, we have

Just kidding, let's use `set!/begin`

To what does this evaluate?

```
(let ([f 0])
  (let ([g 34])
    (begin
     (set! f g)
     f)))
```

A. 0

B. 34

C. An error

To what does this evaluate?

```
(let ([f 0])
  (let ([g (λ (x) (add1 x))])
    (begin
     (set! f g)
     (f 5))))
```

A. 0

B. 1

C. 5

D. 6

E. An error

To what does this evaluate?

```
(let ([f 0])
  (let ([g (λ (x)
             (if (>= x 10) 10 (f (add1 x))))])
    (begin
     (set! f g)
     (f 5))))
```

A. 0

B. 5

C. 10

D. 11

E. An error

Let's draw the environments

```
(let ([f 0])
  (let ([g (λ (x)
             (if (>= x 10)
                 10
                 (f (add1 x))))])
    (begin
     (set! f g)
     (f 5))))
```

Write factorial without letrec

```
(let ([fact 0])
  (let ([placeholder (λ (n)
                       (if (= n 0)
                           1
                           (* n (fact (sub1 n)))))])
    (begin
     (set! fact placeholder)
     (fact 5))))
```

Mutual recursion

```
(letrec ([even? (lambda (x)
                  (cond [(= 0 x) #t]
                        [(= 1 x) #f]
                        [else (odd? (sub1 x))]))]
         [odd? (lambda (x)
                 (cond [(= 0 x) #f]
                       [(= 1 x) #t]
                       [else (even? (sub1 x))]))])
  (odd? 23))
```

Mutual recursion without `letrec`

```
(let ([even? 0]
      [odd? 0])
  (let ([f (lambda (x)
             (cond [(= 0 x) #t]
                   [(= 1 x) #f]
                   [else (odd? (- x 1))]))]
        [g (lambda (x)
             (cond [(= 0 x) #f]
                   [(= 1 x) #t]
                   [else (even? (- x 1))]))])
    (begin
     (set! even? f)
     (set! odd? g)
     (odd? 23))))
```

General transformation

Replace

```
(letrec ([f1 exp1] … [fn expn])
  body)
```

with

```
(let ([f1 0] … [fn 0])
  (let ([g1 exp1] … [gn expn])
    (begin
     (set! f1 g1)
     …
     (set! fn gn)
     body)))
```

General transformation

Replace

```
(letrec ([f1 exp1] … [fn expn])
  body)
```

with

```
(let ([f1 0] … [fn 0])
  (let ([g1 exp1] … [gn expn])
    (begin
     (set! f1 g1)
     …
     (set! fn gn)
     body)))
```

We need some new symbols!

Generating symbols

`(gensym)`

We can use `(gensym)` to generate new, unused symbols

```
> (gensym)
'g75075
> (gensym)
'g75106
```

A common mistake with `gensym`

Every time you call (gensym), you get a new symbol

If you transform `(letrec ([f …]) …)` into

```
(let ([f 0])
  (let ([(gensym) …])
    (begin
     (set! f (gensym))
     …)))
```

your code will fail to work because the two symbols will be different!

Final MiniScheme grammar

EXP →	number	parse into `lit-exp`
|	symbol	parse into `var-exp`
|	(if *EXP EXP EXP*)	parse into `ite-exp`
|	(let (*LET-BINDINGS*) *EXP*)	parse into `let-exp`
|	(letrec (*LET-BINDINGS*) *EXP*)	transform into equivalent `let-exp`
|	(lambda (*PARAMS*) *EXP*)	parse into `lambda-exp`
|	(set! symbol *EXP*)	parse into `set-exp`
|	(begin *EXP**)	parse into `begin-exp`
|	(*EXP EXP**)	parse into `app-exp`

LET-BINDINGS → *LET-BINDING**

LET-BINDING → [symbol *EXP*]*

PARAMS → symbol*

Parsing letrec expressions

`(letrec ([f1 exp1] … [fn expn]) body)`

We have three parts

- `syms = (f1 … fn) = (map first (second input))`
- `exps = (exp1 … expn) = (map second (second input))`
- `body = (third input)`

We need to construct several parts from these

- The outer `let`: `(let ([f1 0] … [fn 0]) …)`
- The inner `let`: `(let ([g1 exp1] … [gn expn]) …)`
- The `set!`s: `(begin (set! f1 g1) … (set! fn gn) …)`

The outer let

`(let ([f1 0] … [fn 0]) …)`

Recall that our `let-exp` has a list of symbols, a list of parsed expressions, and a parsed body

We already got the symbols: `(f1 … fn) = syms`

For the parsed expressions: `(map (λ (s) (lit-exp 0)) syms)`

The parsed body is going to be another `let-exp`

The inner let

`(let ([g1 exp1] … [gn expn]) …)`

For the symbols: `new-syms = (map (λ (s) (gensym)) syms)`

For the parsed expressions: `(map parse exps)`

The parsed body is a begin expression

The begin expression

```
(begin (set! f1 g1) … (set! fn gn) body)
```

Recall that `begin-exp` takes a list of parsed expressions

Three reasonable options

- Generate the `set!`s via `(map (λ (s new-s) …) syms new-syms)`
  Append `(list (parse body))`
- Write your own recursive procedure to build the list
- Use `foldr`

```
(foldr (λ (s new-s acc)
          (cons … acc))
       (list (parse body))
       syms
       new-syms)
```

A (mostly) complete example

```
(letrec ([length (lambda (lst)
                   (if (null? lst)
                       0
                       (add1 (length (cdr lst)))))])
  (length (list 10 20 30)))
```

parses to

```
'(let-exp (length)
          ((lit-exp 0))
          (let-exp (g75784)
                   ((lambda-exp (lst) (ite-exp …)))
                   (begin-exp
                    ((set-exp length (var-exp g75784))
                     (app-exp (var-exp length) (…))))))
```

And that's it!

We don't need to change `eval-exp` at all because we already know how to evaluate let-, set-, and begin-expressions.