

# CS 241: Systems Programming

## Lecture 27. System Calls II

Spring 2020

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# Creating a new process

Two schools of thought

- ▶ Windows way: single system call
  - `CreateProcess("calc.exe", /* other params */)`
- ▶ Unix way: two (or more) system calls
  - Create a copy of the currently running process: `fork()`
  - Transform the copy into a new process:  
`execve("/usr/bin/bc", args, env)`

# Process IDs

Every Unix process has a unique identifier

- ▶ Integer, used to index into a kernel process table
- ▶ `$ ps ax # Print a list of all running processes and their PIDs`

```
pid_t getpid(void);
```

Every process has a parent process

- ▶ processes are "reparented" to the `init` process if your parent already died

```
pid_t getppid(void);
```

# Running another program

```
int execve(char const *path, char *const argv[],  
           char *const envp[]);
```

- ▶ Last element of `argv[ ]` and `envp[ ]` must be 0 (**NULL**)
- ▶ If successful, `execve` won't return, instead, the OS will remove all of the process's code and data and load the program from `path` in its place and start running that
- ▶ The PID of the process doesn't change
- ▶ The open file descriptors remain open (unless marked close on exec)
- ▶ Returns `-1` and sets **errno** on error

# The types of argv and envp

`execve(path, argv, envp)` does not modify its arguments

For historical reasons, `argv` and `envp` have type

- ▶ `char *const[ ]` — this is a constant pointer to `char *`
- ▶ We really want `char const *const[ ]` which is a constant pointer to `char const *`
- ▶ Normally, we pass a `char *argv[ ]` array (no `const`)

# The types of argv and envp

We can deal with this in one of two ways

- ▶ For historical reasons, we can assign string literals to **char \***  
**char** \*s = "foo"; // normally char const \*s = "foo";
- ▶ We can **cast** a **char const \*** to a **char \***  
**char const** \*t = /\* ... \*/;  
**char** \*u = (**char** \*)t;
- ▶ If you omit the cast, you get a compiler warning; compiler warnings should not be ignored

```

#include <err.h>
#include <stdlib.h>
#include <unistd.h>

void run_with_args(char const *program) {
    char *args[] = {
        (char *)program,           // argv[0]
        "This is one argument",    // argv[1]
        "two",                     // argv[2]
        "three",                   // argv[3]
        0,                         // argv[4] is NULL, end of args
    };
    char *env[] = { 0 }; // Empty environment.
    execve(program, args, env);
    err(EXIT_FAILURE, "%s", args[0]);
}

int main(int argc, char *argv[]) {
    run_with_args(argc == 1 ? "/bin/echo" : argv[1]);
}

```

# exec(3) family

```
int execl(const char *path, const char *arg0, ...,  
          (char *)0);
```

```
int execl_e(const char *path, const char *arg0, ...,  
            (char *)0, char *const envp[]);
```

```
int execl_p(const char *program, const char *arg0, ...,  
            (char *)0);
```

```
int execv(const char *path, char *const argv[]);
```

```
int execvp(const char *program, char *const argv[]);
```

- ▶ execl, execl\_e, execl\_p take 0-terminated variable number of arguments
- ▶ The argv and envp arrays must be 0-terminated
- ▶ execl\_p and execvp search PATH for the program
- ▶ glibc has an execvpe which is like execve but searches the PATH



Which of the following statements about `execve ( )` is false?

- A. If `execve()` is successful, the new program replaces the calling program.
- B. The file descriptors that were open before `execve()` are open in the new program (except for those marked as close on exec).
- C. If `execve ( )` has an error, it returns -1 and sets **`errno`**.
- D. If `execve ( )` is successful, it returns 0.

# Creating a new process

```
#include <unistd.h>
#include <sys/types.h>
```

```
pid_t fork(void);
```

Creates an (almost) identical copy of the running program with one big exception

- ▶ Returns 0 to the child but PID of child to the parent
- ▶ -1 on error and sets **errno**

This includes a copy of memory, code, file descriptors and most other bit of process state (but not all)

```

#include <sys/types.h>
#include <sys/wait.h>
#include <err.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

void whoami(char const *str) {
    pid_t self = getpid();
    pid_t parent = getppid();
    printf("%s: pid=%d ppid=%d\n",
        str, self, parent);
}

```

```

int main(void) {
    whoami("Prefork");
    pid_t pid = fork();
    if (pid < 0)
        err(EXIT_FAILURE, "fork");
    if (pid == 0) {
        whoami("Child");
    } else {
        whoami("Parent");
        int status;
        wait(&status);
    }
    return 0;
}

```

```
#include <sys/types.h>
```

```
Prefork: pid=48627 ppid=28834
```

```
Parent: pid=48627 ppid=28834
```

```
Child: pid=48628 ppid=48627
```

```
#include <stdlib.h>
```

```
#include <unistd.h>
```

```
void whoami(char const *str) {  
    pid_t self = getpid();  
    pid_t parent = getppid();  
    printf("%s: pid=%d ppid=%d\n",  
        str, self, parent);  
}
```

```
int main(void) {  
    whoami("Prefork");  
    pid_t pid = fork();  
    if (pid < 0)  
        err(EXIT_FAILURE, "fork");  
    if (pid == 0) {  
        whoami("Child");  
    } else {  
        whoami("Parent");  
        int status;  
        wait(&status);  
    }  
    return 0;  
}
```

# fork/exec

Usually used together

`fork` to create a duplicate process

`exec` (one of the `exec` family that is) to run a new program

`fork` and `exec` both preserve file descriptors

- ▶ This is how `bash` operates: it forks, sets file descriptors, and execs

After a `fork`, you have two copies of a program, the parent and the child, and...

- A. Either the parent or the child must call `exec ( )` immediately
- B. The parent gets a PID and the child gets a 0 as return values
- C. The child gets a PID and the parent gets a 0 as return values
- D. Both parent and child get PIDs as the return values
- E. Both parent and child must call `exec` to proceed

# Process exit status

Can wait for a child process to die (or be stopped, e.g., by a debugger)

```
#include <sys/wait.h>
```

```
int status;
```

```
pid_t pid = wait(&status);
```

Suspends execution until child terminates, returns the PID of the child

# Checking exit status

Use macros to examine exit status

**WIFEXITED**( status )

- ▶ True if the process terminated normally

**WEXITSTATUS**( status )

- ▶ Returns actual return/exit value if **WIFEXITED**( status ) is true

**WIFSIGNALED**( status )

- ▶ True if the process was terminated by a signal (e.g., **SIGINT** from ctrl-C)

**WTERMSIG**( status )

- ▶ Returns the signal that terminated the process if **WIFSIGNALED**( status )



# strace(1)

strace is a Linux program that prints out the system calls a program uses

- ▶ `-e trace=open,openat,close,read,write` will trace those system calls
- ▶ `-f` will trace children too
- ▶ `-s size` will show up to `size` bytes of strings

```
$ strace -e trace=open,openat,close,read,write cat Makefile
...
openat(AT_FDCWD, "Makefile", O_RDONLY) = 3
read(3, "CC := clang\nCFLAGS := -Wall -std"..., 1048576) = 176
write(1, "CC := clang\nCFLAGS := -Wall -std"..., 176) = 176
read(3, "", 1048576) = 0
close(3) = 0
...
```

# In-class exercise

<https://checkoway.net/teaching/cs241/2020-spring/exercises/Lecture-27.html>

Grab a laptop and a partner and try to get as much of that done as you can!